

Finding ECM friendly curves: A Galois approach

Sudarshan SHINDE

Institut de Mathématiques de Jussieu - Paris Rive Gauche

Motivation - ECM algorithm

Algorithm 1 ECM (H. Lenstra 1985)

INPUT : n with at least two different prime factors

OUTPUT : a non-trivial factor of n .

- 1: $B \leftarrow B_n, m \leftarrow B!$
 - 2: **while** No factor is found **do**
 - 3: $P \leftarrow (x, y) \in \mathbb{Z}/n\mathbb{Z} \times \mathbb{Z}/n\mathbb{Z}$ and an elliptic curve E on $\mathbb{Z}/n\mathbb{Z}$
 such that and $P \in E(\mathbb{Z}/n\mathbb{Z})$.
 - 4: $P_m \leftarrow [m]P = (x_m : y_m : z_m) \bmod n$
 - 5: $g \leftarrow \gcd(z_m, n)$
 - 6: **if** $g \notin \{1, n\}$ **then return** g
 - 7: **end if**
 - 8: **end while**
-

Correctness

Idea

Let p be an unknown prime factor of n . If $\text{ord}(P)$ in $E(\mathbb{F}_p)$ divides $B!$, then

$$(x_B : y_B : z_B) \equiv (0 : 1 : 0) \pmod{p}.$$

In this case p divides $\gcd(z_B, n)$.

Sufficient condition

$\#E(\mathbb{F}_p)$ is B -smooth i.e. all its prime factors are $< B$.

Theorem (Hasse)

Let E be an elliptic curve and p be a prime. Then,

$$p + 1 - 2\sqrt{p} \leq \#E(\mathbb{F}_p) \leq p + 1 + 2\sqrt{p}.$$

Probability of success

- ECM succeeds if $\#E(\mathbb{F}_p)$ is B -smooth.
- $\#E(\mathbb{F}_p) \approx p$ (Hasse)

Theorem (Lenstra)

For $p > 3$, let $S_p = \{s : |s - (p + 1)| \leq \sqrt{p} \text{ and } s \text{ is } B\text{-smooth}\}$.
Then the probability ECM succeeds is at least $\frac{c}{\log(p)} \frac{\#S-2}{2\sqrt{p+1}}$.

In other words, probability that a particular curve succeeds is comparable with the probability of finding a B -smooth number in the interval $(p + 1 - \sqrt{p}, p + 1 + \sqrt{p})$.

Improved ECM algorithm

Algorithm 2 Practical version of ECM

INPUT : n with at least two different prime factors

OUTPUT : a non-trivial factor of n .

- 1: $B \leftarrow B_n, m \leftarrow B!$
 - 2: **while** No factor is found **do**
 - 3: $E/\mathbb{Q} \leftarrow$ an elliptic curve **from a family** and $P \in E(\mathbb{Q})$.
 - 4: $P_m \leftarrow [m]P = (x_m : y_m : z_m) \bmod n$
 - 5: $g \leftarrow \gcd(z_m, n)$
 - 6: **if** $g \notin \{1, n\}$ **then return** g
 - 7: **end if**
 - 8: **end while**
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Idea of Montgomery

Question : What if $\#E(\mathbb{F}_p)$ is even for all primes p ?

Theorem : If m divides torsion order of $E(\mathbb{Q})$ then m divides $\#E(\mathbb{F}_p)$ for almost all p .

Montgomery heuristic

Definition

Let E be an elliptic curve, ℓ be a prime and B be a sufficiently large integer. We define empirical average valuation,

$$\bar{v}_\ell(E) = \frac{\sum_{p < B} (\text{val}_\ell(\#E(\mathbb{F}_p)))}{\#\{p < B\}}$$

Heuristic

Curves with larger average valuation are ECM-friendly.

How to improve average valuation ?

Some ways

- 1 Montgomery (1985), Suyama (1985), Atkin et Morain (1993), Bernstein et al (2010) : Torsion points over \mathbb{Q}

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- ② Brier and Clavier (2010) : Torsion points over $\mathbb{Q}(i)$

$$\bar{v}_2(\#E(\mathbb{F}_p)) = \frac{1}{2}\bar{v}_2(\#E(\mathbb{F}_p) | p \equiv 1 \pmod{4}) + \frac{1}{2}\bar{v}_2(\#E(\mathbb{F}_p) | p \equiv 3 \pmod{4})$$

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- ③ Barbulescu et al (2012) : Better average valuation without additional torsion points by reducing the size of a "specific" Galois group.

Preliminaries

Definition (m -torsion field)

Let E be an elliptic curve on \mathbb{Q} , m a positive integer. The m -torsion field $\mathbb{Q}(E[m])$ is defined as the smallest extension of \mathbb{Q} containing all the m -torsion points.

As $E(\bar{\mathbb{Q}})[m] \simeq \mathbb{Z}/m\mathbb{Z} \times \mathbb{Z}/m\mathbb{Z}$, $G = \text{Gal}(\mathbb{Q}(E[m])/\mathbb{Q})$ is always a subgroup of $\text{Aut}(\mathbb{Z}/m\mathbb{Z} \times \mathbb{Z}/m\mathbb{Z}) = \text{GL}_2(\mathbb{Z}/m\mathbb{Z})$.

Theorem (Serre)

Let E be an elliptic curve without complex multiplication.

- For all primes ℓ outside a finite set depending on E and for all $k \geq 1$, $\text{Gal}(\mathbb{Q}(E[\ell^k])/\mathbb{Q}) = \text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z})$
- For all primes ℓ and $k \geq 1$, the index $[\text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z}) : \text{Gal}(\mathbb{Q}(E[\ell^k])/\mathbb{Q})]$ is non-decreasing and bounded by a constant depending on E and ℓ .

How to improve average valuation ?

Theorem (Barbulescu et al. 2012)

Let ℓ be a prime and E_1 and E_2 be two elliptic curves. If $\forall n \in \mathbb{N}, \text{Gal}(\mathbb{Q}(E_1[\ell^n])) \simeq \text{Gal}(\mathbb{Q}(E_2[\ell^n]))$ then $\bar{v}_\ell(E_1) = \bar{v}_\ell(E_2)$.

Thus in order to change the average valuation,
we must change $\text{Gal}(\mathbb{Q}(E[\ell^n]))$ for at least one n .

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Exemple

Family	Torsion	\bar{v}_2	Primes found between $2^{15}, 2^{22}$
Suyama	$\mathbb{Z}/6\mathbb{Z}$	$10/3$	4069
Suyama - 11	$\mathbb{Z}/6\mathbb{Z}$	$11/3$	4756 (16% more)

Suyama-11 is implemented in GMP-ECM.

Constructing torsion field - Division polynomials

Definition - Theorem

For an elliptic curve E and a an integer m , we define the m -division polynomial as

$$\Psi_{(E,m)}(X) = \prod_{(x_P, \pm y_P) \in E[m] - O} (X - x_P) \in \mathbb{Q}[X],$$

and the exact m -division polynomial as

$$\Psi_{(E,m)}^{\text{exact}}(X) = \prod_{(x_P, \pm y_P) \text{ of order } m} (X - x_P) \in \mathbb{Q}[X].$$

We have $\deg(\Psi_{(E,m)}) = \frac{m^2+2-3\eta}{2}$ where $\eta = m \bmod 2$.

Example

Let $E : y^2 = x^3 + ax + b$ then $\Psi_{(E,3)} = x^4 + 2ax^2 + 4bx - \frac{1}{3}a^2$

Constructing prime-power torsion field

Given $E : y^2 = x^3 + ax + b$ and a prime-power ℓ^n , we construct $\mathbb{Q}(E[\ell^n])$ recursively :

Constructing $\mathbb{Q}(E[\ell])$

$$\mathbb{Q} \hookrightarrow \mathbb{Q}(x_1) \hookrightarrow \mathbb{Q}(x_1, x_2) \hookrightarrow \mathbb{Q}(x_1, x_2, y_1) \hookrightarrow \mathbb{Q}(x_1, x_2, y_1, y_2) = \mathbb{Q}(E[\ell])$$

where the polynomials defining the extensions are ;

- 1 (An irreducible factor of) $\Psi_{(E,\ell)}$
- 2 An irreducible factor of $\Psi_{(E,\ell)}$ on $\mathbb{Q}(x_1)$.
- 3 $f_1(y) = y^2 - (x_1^3 + ax_1 + b)$.
- 4 $f_2(y) = y^2 - (x_2^3 + ax_2 + b)$

Once we have $\mathbb{Q}(E[\ell^{n-1}])$, we construct $\mathbb{Q}(E[\ell^n])$ by the same method using $\Psi_{(E,\ell^n)}^{\text{exact}}$ over $\mathbb{Q}(E[\ell^{n-1}])$.

Inverse Galois problem - Main theorem

Definition (Resolvent polynomial)

Let G be a subgroup of \mathcal{S}_n and $F(X_1, \dots, X_n) \in K[X_1, \dots, X_n]$ such that $G = \{\sigma \in \mathcal{S}_n \mid F(X_{\sigma(1)}, \dots, X_{\sigma(n)}) = F(X_1, \dots, X_n)\}$. For a polynomial P , we define the resolvent polynomial

$$R_G(F, P)(X) = \prod_{\sigma \in \mathcal{S}_n/G} (X - F(\theta_{\sigma(1)}, \dots, \theta_{\sigma(n)})),$$

where $\theta_1, \dots, \theta_n$ are the roots of P in \bar{K}

Theorem

Let P, G, F be as above. Then,

- 1 $R_G(F, P)(X) \in K[X]$.
- 2 If $\text{Gal}(P) \subset G$ then $R_G(F, P)(X)$ has a root in K and if $R_G(F, P)(X)$ has a *simple* root in K then $\text{Gal}(P) \subset G$ upto conjugacy.

The theorem over $K = \mathbb{Q}(a_1, \dots, a_n) \Rightarrow$ inverse Galois problem.

The particular case of division polynomials

$$\text{Split}(\Psi_{(E,m)}) \subset \mathbb{Q}(E[m])$$

Particular case

- $m = 2^2$: **Theorem** : For a Montgomery curve $(By^2 = X^3 + Ax^2 + x)$, $\text{Gal}(\Psi_4) \neq \mathbb{Z}/4\mathbb{Z}$.
- $m = 3$ **Theorem** : For any curve, if Ψ_3 is irreducible and $\text{Gal}(\mathbb{Q}(E[3]/\mathbb{Q})) \neq \text{GL}_2(\mathbb{Z}/3\mathbb{Z})$ then $\#\text{Gal}(\Psi_3) = 16$.

Remark

When $P = \Psi_{(E,\ell)}$ of degree $n = \frac{\ell^2-1}{2}$, we have $\deg(R_G) = [S_n : G] > [S_n : \text{GL}_2(\mathbb{Z}/\ell\mathbb{Z})] > \text{exponential}(\ell^2)$

The division polynomial is not a random polynomial.

Algorithmic search of infinite families

Computer algebra approach

Idea : Formal construction of torsion field and sufficient condition that its Galois group is generic.

Sufficient condition : When all the following extensions have generic degrees.

$$\begin{aligned} K_4 &= \mathbb{Q}(a, b)(x_1, x_2, y_1, y_2) = \mathbb{Q}(a, b)(E[\ell]) \\ &\quad \left| P_4 = y^2 - (x_2^3 + ax_2 + b) \right. \\ K_3 &= \mathbb{Q}(a, b)(x_1, x_2, y_1) \\ &\quad \left| P_3 = y^2 - (x_1^3 + ax_1 + b) \right. \\ K_2 &= \mathbb{Q}(a, b)(x_1, x_2) \\ &\quad \left| P_2 = \text{a factor of } \Psi \text{ of degree } \frac{\ell^2 - \ell}{2} \right. \\ K_1 &= \mathbb{Q}(a, b)(x_1) \\ &\quad \left| P_1 = \Psi \text{ of degree } \frac{\ell^2 - 1}{2} \right. \\ K_0 &= \mathbb{Q}(a, b) \end{aligned}$$

Test if the above four polynomials are irreducible.

Non-generic Galois image

Example :

Let $E : y^2 = x^3 + ax + b$ be an elliptic curve. Then $\Psi_3 = x^4 + 2ax^2 + 4bx - \frac{1}{3}a^2$. We consider a partition of 4 of length 2.

- For $[2, 2]$, we write,

$$x^4 + 2ax^2 + 4bx - \frac{1}{3}a^2 = (x^2 + e_2x + e_1)(x^2 + f_2x + f_1)$$

and equate the coefficients on both sides. We get a system of polynomial equations,

$$\begin{cases} e_2 + f_2 = 0 \\ e_2f_2 + e_1 + f_1 = 2a \\ e_1f_2 + e_2f_1 = 4b \\ e_1f_1 = -1/3 a^2 \end{cases} \Rightarrow \begin{cases} f_2 = -e_2 \\ f_1 = 2a + e_2f_2 - e_1 \\ e_1(e_2^2 + 2a - e_1) + \frac{1}{3}a^2 = 0. \\ e_2^6 + 4ae_2^4 + \frac{16}{3}e_2^2a^2 - 16b^2 = 0 \end{cases}$$

Thus, if the polynomial $3x^6 + 12ax^4 + 16a^2x^2 - 48b^2$ does not have a root, then the factorization pattern of Ψ_3 is not $[2, 2]$.

Algorithm

Algorithm 3 Finding families

INPUT : A prime ℓ

OUTPUT : Necessary polynomial conditions in a and b such that $\text{Gal}(\mathbb{Q}(E[\ell]))$ is non-generic for an elliptic curve E over $\mathbb{Q}(a, b)$

- 1: **for** $i \in \{1, 2, 3, 4\}$ **do**
- 2: $F_i \leftarrow$ absolute polynomial of K_{i-1}
- 3: **for** $r \in$ partitions of $\deg(P_i)$ **do**
- 4: $S_{i,r} \leftarrow$ System of polynomial equations in a, b and a root of F_i arising from equating coefficients
- 5: $C_{i,r} \leftarrow$ Triangulation of $S_{i,r}$ (Resultant)
- 6: ▷ Necessary for factorization pattern of P_i to be r .
- 7: **end for**
- 8: **end for**
- 9: **return** Set of $C_{i,r}$

A unified presentation

Montgomery (1992) : "The table entries were found in an **ad hoc** manner, so I make no claim completeness."

Kruppa (2007) : "The choice of $\sigma = 11$, which **surprisingly** leads to higher average exponent.."

Barbulescu et al (2012) : ".. **suggests** that by imposing equations on the parameters a and d we can improve the torsion properties."

"..**By trying** to force one of these three polynomials to split, we found four families."

Case $l = 3$

Theorem

Let $E : y^2 = x^3 + ax + b$ be a rational elliptic curve with $ab \neq 0$. Let Ψ_3 be its 3-division polynomial and Δ its discriminant. Then we have,

Fact. Pattern of Ψ_3	Condition(s)	index
(1, 1, 2)	C_1 and a 3-torsion point	24
(1, 1, 2)	C_1	12
(1, 3)	$C_{2'}$ or [C_2 and a 3-torsion point]	8
(1, 3)	C_2	4
(2, 2)	C_3	6
(4)	C_4	3

$$C_1 = 27x^{12} + 594ax^{10} + 972bx^9 + 4761a^2x^8 + 14256abx^7 + \dots + 324ab(587a^3 + 3456b^2)x - 5329a^6 + 162432b^2a^3 + 1492992b^4$$

$$C_{2'} = x^{16} - 24bx^{12} + 6\Delta x^8 - 3\Delta^2$$

$$C_2 = 3x^4 + 6ax^2 + 12bx - a^2$$

$$C_3 = 3x^6 + 12ax^4 + 16a^2x^2 - 48b^2$$

$$C_4 = x^3 - 2\Delta \text{ i.e. the } j \text{ of } E \text{ is a cube.}$$

From conditions to families of curves

Remark

For every case, we got the equations of type $\exists x \in \mathbb{Q}$ such that $C(a, b, x) = 0$.

From surface to curve

$(a, b) \sim (as^4, bs^6)$ (Same elliptic curve over \mathbb{Q}), So essentially $C(a, b, x)$ is a plane curve. Replacing a and b by $a(j)$ and $b(j)$ or by random linear polynomials in t , we obtain a curve $C(a(t), b(t), x)$. This curve describes infinitely many elliptic curves having the same Galois image.

Classical results on curves

Let C be a non-singular curve of genus g . Then if,

- $g \geq 2$, C has finitely many points. [Faltings]
- $g = 0$, if there is a point, there are infinitely many.
- $g = 1$, if there is a point, C can be put in Weierstrass form with rank r .

From conditions to families of curves : Example

Let $E : y^2 = x^3 + ax + b$ be an elliptic curve. We saw that if Ψ_3 factors into two quadratic factors then

$C = 3x^6 + 12ax^4 + 16a^2x^2 - 48b^2$ has a root.

If we put $b = 2a$, we get $C = 3x^6 + 12ax^4 + 16a^2x^2 - 192a^2$.

This curve is of genus 0. We get a parametrization

$$a(t) = \frac{27t^3(19t+2)^3}{(242t^2+54t+3)(271t^2+57t+3)^2} \text{ and } b(t) = 2a(t).$$

Computing the generic valuation of a family

Theorem

Let $E_t : y^2 = x^3 + a(t)x + b(t)$ such that $\text{Gal}(\mathbb{Q}(t)(E_t[\ell])/\mathbb{Q}(t)) \subseteq H$. If $\exists t_0 \in \mathbb{Q}$ such that $\#\text{Gal}(\mathbb{Q}(E_{t_0}[\ell])/\mathbb{Q}) = \#H$ then $\text{Gal}(\mathbb{Q}(t)(E_t[\ell])/\mathbb{Q}(t)) = H$.

Proof

Let $E_t : y^2 = x^3 + a(t)x + b(t)$ and $p = t - t_0$.

$$\begin{array}{ccc}
 \text{Gal}(\mathbb{Q}(t)(E_t[\ell])/\mathbb{Q}(t)) & \xleftarrow{\cong} \text{Dec}(\mathfrak{p}) & \xrightarrow{\text{eval}_{t=t_0}} \text{Gal}(\mathbb{Q}(E_{t_0}[\ell])/\mathbb{Q}) \\
 \downarrow \rho & & \downarrow \rho \\
 \text{GL}_2(\mathbb{Z}/\ell\mathbb{Z}) & \xrightarrow{\cong} & \text{GL}_2(\mathbb{Z}/\ell\mathbb{Z})
 \end{array}$$

where Dec is defined below.

Let K be a Galois extension of $\mathbb{Q}(t)$. Let $\mathfrak{p} \in \mathbb{Q}(t)$ and \mathfrak{p} be an ideal of K above \mathfrak{p} . We define $\text{Dec}(\mathfrak{p}) = \{\sigma \in \text{Gal}(K/\mathbb{Q}) \mid \sigma(\mathfrak{p}) = \mathfrak{p}\}$.

Valuations for $\ell = 3$

We obtain $g = 0$ for all the families for $\ell = 3$.

Theorem

Let $E : y^2 = x^3 + ax + b$, $ab \neq 0$ be a rational elliptic curve. Then the generic average valuation $\bar{v}_3(E)$ is $87/128 \approx 0.68$, except when one the following cases occurs.

A parametrization	Example (a, b)	Valuation
a, b complicated.	$(5805, -285714)$	$33/16 \approx 2.06$
a, b complicated.	$(284445, 97999902)$	$45/32 \approx 1.41$
$a = 3t^2, b = -243t^6 + 162t^4 - 9t^2/36$	$(3, -11)$	$27/16 \approx 1.69$
$a = \frac{-192t^3 - 254803968}{t^4}, b = \frac{-t^6 - 5308416t^3 - 4696546738176}{3t^6}$	$(-254804160, -\frac{4696552046593}{3})$	$27/16 \approx 1.69$
$a = \frac{-36t(t+2)^3}{(t^2+4t+1)^2}, b = 2a$	$(\frac{-4608}{169}, \frac{-9216}{169})$	$39/32 = 1.22$
$a = \frac{27t^3(19t+2)^3}{(242t^2+54t+3)(271t^2+57t+3)^2}, b = 2a$	$(\frac{250047}{32758739}, \frac{500094}{32758739})$	$69/64 \approx 1.08$
$a = \frac{216}{(t^3-8)}, b = 2a$	$(\frac{-216}{7}, \frac{-432}{7})$	$69/128 \approx 0.54$

Cryptographic application

Popular parametrizations

- Montgomery $By^2 = x^3 + Ax^2 + x$ or $y^2 = x^3 + \frac{3-A^2}{3B^2}x + \frac{2A^3-3A}{27B^3}$
- Edwards $ax^2 + y^2 = 1 + dx^2y^2$ or $y^2 = x^3 + \frac{3-\alpha^2}{3\beta^2}x + \frac{2\alpha^3-3\alpha}{27\beta^3}$
where $\alpha = -2\frac{a+d}{a-d}$ and $\beta = \frac{4}{a-d}$.
- Hessian $y^2 + axy + by = x^3$ or $y^2 = x^3 + (-27a^4 + 648ab)x + (54a^6 - 1944a^3b + 11664b^2)$.

Goal

- **INPUT** : A number field K , a prime ℓ and $a(\alpha, \beta)$ and $b(\alpha, \beta)$.
- **OUTPUT** : Complete list of equations of negligible density necessary for non-generic valuation $\bar{v}_\ell(E_{a,b})$.

Valuation $m = 4$, Montgomery curve

Theorem

Let $E : By^2 = x^3 + Ax^2 + x$ be a rational elliptic curve with $B(A^2 - 4) \neq 0$. Then the generic average valuation $\bar{v}_2(E)$ is $^{10}/_3 \approx 3.33$, except,

- If $A^2 - 4 \neq \square$ i.e. $E(\mathbb{Q})[2] \neq \mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/2\mathbb{Z}$, we note Ψ be the quartic factor of its 4-division polynomial. Then we have,

Fact. Pat. of Ψ	Condition(s)	Index	Valuation
(2, 2)	$A = -2 \frac{t^4 - 4}{t^4 + 4}$	24	$^{10}/_3 \approx 3.33$
(4)	$\frac{A \pm 2}{B} = \pm \square$ or $\frac{4B^2}{A^2 - 4} = -t^4$	12	$^{11}/_3 \approx 3.67$

- If $A^2 - 4 = \square$ i.e. if $A = \frac{t^2 + 4}{2t}$. Then we have,

Fact. Pat. of Ψ	Condition(s)	Index	Valuation
(1, 1, 2)	$A = \frac{t^4 + 24t^2 + 16}{4(t^2 + 4)t}$ and $B = -t(t^2 + 4)\square$	48	$^{14}/_3 \approx 4.67$
(1, 1, 2)	$A = \frac{t^4 + 24t^2 + 16}{4(t^2 + 4)t}$	24	$^{23}/_6 \approx 3.83$
(2, 2)	$A = \frac{t^2 + 4}{2t}$ and $\frac{A \pm 2}{B} = \square$	24	$^{13}/_3 \approx 4.33$
(2, 2)	$A = \frac{t^2 + 4}{2t}$	12	$^{11}/_3 \approx 3.67$

Modular forms approach

Theorem (Sutherland, Zywina)

Let E be an elliptic curve and $H \subset GL_2(\mathbb{Z}/\ell^n\mathbb{Z})$ such that $-1 \in H$. Then there exists a polynomial $X_H(j, t)$ such that $\text{Gal}(\mathbb{Q}(E[\ell^n])/\mathbb{Q}) \subset H$ if and only if $\exists t_0 \in \mathbb{Q}$ such that $X_H(j(E), t_0) = 0$.

Fast computations of X_H

[1] Jeremy Rouse and David Zureick-Brown, "Elliptic curves over \mathbb{Q} and 2-adic images of Galois" (2015)

- Complete description of possible 2-adic Galois images.

[2] Andrew Sutherland and David Zywina, "Modular curves of prime-power level with infinitely many rational points" (2017)

- Complete description of possible ℓ -adic Galois images contained in subgroups containing -1 .

Example

Curve	$j(E)$	$\#\text{Gal}(\mathbb{Q}(E[3])/\mathbb{Q})$	\bar{v}_3
$y^2 = x^3 - 336x + 448$	1792	12	$39/32$
$y^2 = x^3 - 7^2 \cdot 336x + 7^3 \cdot 448$	1792	6	$54/32$

The modular forms approach does not work for arbitrary H .

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$y^2 = x^3 - 7^2 \cdot 336x + 7^3 \cdot 448$	1792	6	$54/32$

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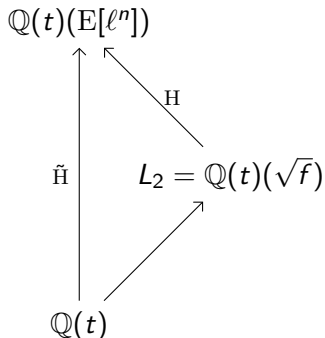
Let H be a subgroup of $\text{GL}_2(\mathbb{Z}/\ell^n\mathbb{Z})$.

	$-1 \notin H$	$-1 \in H$
$l = 2$	[1]	[1], [2]
$l \neq 2$		[2]

Our contribution

Complete list of elliptic curves having non-generic Galois image not containing -1 .

Let \tilde{H} be subgroup of $GL_2(\mathbb{Z}/\ell^n\mathbb{Z})$ containing -1 and let H be subgroup of \tilde{H} such that $\tilde{H} = \langle H, -1 \rangle$.



Cryptographic applications

A criterion to choose curves

Notation : $s \sim t$ if $t - \sqrt{t} < s < t + \sqrt{t}$.

p is fixed and E varies (H. Lenstra)

$$\text{Prob}(\#E(\mathbb{F}_p) \text{ is } B\text{-smooth}) = \frac{1}{\mathcal{O}(\log p)} \text{Prob}(a \text{ random integer } \sim p \text{ is } B\text{-smooth}).$$

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E fixed and p varies in $[n - \sqrt{n}, n + \sqrt{n}]$

Can we claim the following ?

$$\text{Prob}(\#E(\mathbb{F}_p) \text{ is } B\text{-smooth}) = \text{Prob}(a \text{ random integer } \sim ne^\alpha \text{ is } B\text{-smooth}).$$

A criterion to choose curves

Notation : $s \sim t$ if $t - \sqrt{t} < s < t + \sqrt{t}$.

p is fixed and E varies (H. Lenstra)

$$\text{Prob}(\#E(\mathbb{F}_p) \text{ is } B\text{-smooth}) = \frac{1}{\mathcal{O}(\log p)} \text{Prob}(a \text{ random integer } \sim p \text{ is } B\text{-smooth}).$$

E fixed and p varies in $[n - \sqrt{n}, n + \sqrt{n}]$

Can we claim the following ?

$$\text{Prob}(\#E(\mathbb{F}_p) \text{ is } B\text{-smooth}) = \text{Prob}(a \text{ random integer } \sim ne^\alpha \text{ is } B\text{-smooth}).$$

Definition

For E an elliptic curve and n, B two integers, $\alpha(E, n, B) \in \mathbb{R}$ is such that

$$\frac{\#\{p \sim n \mid \#E(\mathbb{F}_p) \text{ is } B\text{-smooth}\}}{\#\{p \mid p \sim n\}} = \frac{\#\{x \sim ne^{\alpha(E, n, B)} \mid x \text{ is } B\text{-smooth}\}}{\#\{x \mid x \sim ne^{\alpha(E, n, B)}\}}.$$

Example

Let $E : y^2 = x^3 + 3x + 5$ and $n = 2^{25}$. We compute α for usual values of B .

$\alpha(E, n, 30)$	-0.79
$\alpha(E, n, 60)$	-0.83
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Theorem (Barbulescu and Lachand (2016))

Let f be a quadratic homogeneous polynomial with certain properties.

$$\text{Prob}(f(n), n \sim N \text{ is } B\text{-smooth}) = \text{Prob}(n \text{ of size } Ne^{\alpha(f)} \text{ is } B\text{-smooth}).$$

Question : Can we make α independent of n and B ?

Definition

Let n and B' be integers. We define B' sifted part of n ,

$$C_{B'}(n) = \frac{n}{\prod_{p \leq B'} p^{v_p(n)}}.$$

In order to render α independent of n and B , we let $n, B \rightarrow \infty$ and replace proportions by the density of Chebotarev. Let $B' < B$.

Montgomery heuristic

If $C_{B'}(x)$ and $C_{B'}(\#E(\mathbb{F}_p))$ are of the same size then x and $\#E(\mathbb{F}_p)$ have the same chances of being B -smooth.

Thus when B' and $x \rightarrow \infty$, we expect,

$$\alpha + \log(n) - \sum_{\ell} \bar{v}_{\ell}(x) \log(\ell) = \log(n) - \sum_{\ell} \bar{v}_{\ell}(E) \log(\ell).$$

This prompts us to define the following.

Formal definition of $\alpha(E)$

Assuming the convergence for now,

Definition

Let E be an elliptic curve and ℓ a prime. Let $\alpha_\ell(E) = \left(\frac{1}{\ell-1} - \bar{v}_\ell(E)\right) \log \ell$. We define,

$$\alpha(E) = \sum_{\ell} \alpha_\ell(E).$$

Existence and computation of $\alpha(E)$

Calculations of $\bar{v}_\ell(E)$ can be done explicitly using the image of ℓ^n -torsion field.

Generic case

Let E_g be such that for all primes ℓ and for all $k \geq 1$, we have

$\text{Gal}(\mathbb{Q}(E_g[\ell^k])/\mathbb{Q}) = \text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z})$. In this case, we get $\bar{v}_\ell(E_g) = \frac{(\ell^3 + \ell^2 - 2\ell - 1)\ell}{(\ell + 1)^2(\ell - 1)^3}$ and numerically $\alpha(E_g) \approx -0.811997734$.

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Non-generic cases

According to a theorem of Serre, for every elliptic curve without complex multiplication, there are only finitely many primes ℓ for which $\text{Gal}(\mathbb{Q}(E[\ell^k])/\mathbb{Q})$ can be different than $\text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z})$. Thus in this case, α differs by only finitely many terms in its defining series.

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Remark

Serre also conjectured that for every prime $\ell > 37$, $\text{Gal}(\mathbb{Q}(E[\ell^k])/\mathbb{Q}) = \text{GL}_2(\mathbb{Z}/\ell^k\mathbb{Z})$. This enables us to compute α for any given curve effectively assuming the conjecture. Andrew Sutherland has verified this conjecture with the curves in Cremona database.

α : An efficient tool

- ① Curves with torsion $\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/8\mathbb{Z}$: For these curves \bar{v}_2 changes from $\frac{14}{9}$ to $\frac{16}{3}$. Thus,

$$\alpha_{\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/8\mathbb{Z}} = \alpha_{generic} + (14/9 - 16/3) \log(2) \approx -3.4355.$$

- ② Suyama-11 family : For these curves, \bar{v}_2 changes from $\frac{14}{9}$ to $\frac{11}{3}$ and \bar{v}_3 changes from $\frac{87}{128}$ to $\frac{27}{16}$. Thus,

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Numerical experiments with α . ($n = 2^{25}$)

- ① Curves with torsion $\mathbb{Z}/2\mathbb{Z} \times \mathbb{Z}/8\mathbb{Z}$.

	n	ne^α	$\#E(\mathbb{F}_p)$	error_n	error_{ne^α}
$B_1 = 30$	0.000518	0.005753	0.005126	889 %	10.89 %
$B_2 = 100$	0.008892	0.03883	0.042573	378.8 %	9.63 %

- ② Suyama-11

	n	ne^α	$\#E(\mathbb{F}_p)$	error_n	error_{ne^α}
$B_1 = 30$	0.000518	0.005133	0.005743	1008 %	11.89 %
$B_2 = 100$	0.008892	0.04013	0.04101	361%	2.19%

Some other families

$j(t)$	$\alpha(E_{j(t)})$
$\frac{(t^9+9t^6+27t^3+3)^3(t^3+3)^3}{(t^6+9t^3+27)t^3}$	-1.5873
$\frac{256(t^8+8t^6+20t^4+16t^2+1)^3}{(t^2+4)(t^2+2)^2t^2}$	-2.2176
$\frac{(t^8-16t^4+16)^3}{(t^4-16)t^4}$	-2.3908
$\frac{-16(t^{16}-16t^8+16)^3}{(t^8-1)t^{32}}$	-2.4486
$\frac{(t^{16}-8t^{14}+12t^{12}+8t^{10}+230t^8+8t^6+12t^4-8t^2+1)^3}{(t^4-6t^2+1)^2(t^2+1)^4(t^2-1)^8t^8}$	-2.6219
$\frac{(t^{16}-8t^{14}+12t^{12}+8t^{10}+230t^8+8t^6+12t^4-8t^2+1)^3}{(t^4-6t^2+1)^2(t^2+1)^4(t^2-1)^8t^8}$	-3.4355

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There are only finitely many values of $\alpha(E)$. And the best among them is approximately -3.43.

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Thank you !