

# Universal elliptic Gauss sums and applications

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# Classical Gauss sum

Let  $q \neq 2$  be a prime,  $\chi : (\mathbb{Z}/q\mathbb{Z})^* \rightarrow \mu_n, n \mid q-1$ ,  $\xi$  an  $n$ -th root of unity and  $\zeta$  a  $q$ -th root of unity,  $\langle g \rangle = \mathbb{F}_q^*$ . A (cyclotomic) Gauss sum is defined as

$$\sum_{i=1}^{q-1} \chi(g^i) \zeta^{g^i} = \sum_{i=1}^{q-1} \xi^{mi} \zeta^{g^i}$$

# Elliptic Curves

- Recall: Elliptic curve  $E$  over finite field  $\mathbb{F}_p$  ( $p \neq 2, 3$ ):  
 $Y^2 = X^3 + AX + B$ .

Identify  $E$  with set of points  $(X, Y) \in \overline{\mathbb{F}}_p \times \overline{\mathbb{F}}_p$  satisfying the equation together with  $\mathcal{O}$ .

- We wish to determine  
 $\#E(\mathbb{F}_p) = \#\{(X, Y) \in \mathbb{F}_p \times \mathbb{F}_p \mid (X, Y) \text{ lies on } E\} \cup \mathcal{O}$ .
- Important problem related to ECC.

# Definitions and Facts

- *l-torsion*:  $E[\ell] = \{P \in E \mid [\ell]P = \mathcal{O}\}$ .  
Later on,  $\ell$  will be prime,  $\ell \neq p$ . In this case

$$E[\ell] \cong \frac{\mathbb{Z}}{\ell\mathbb{Z}} \times \frac{\mathbb{Z}}{\ell\mathbb{Z}}.$$

- *Frobenius endomorphism*:

$$\phi_p : E \rightarrow E, \quad (X, Y) \mapsto (X^p, Y^p)$$

By restriction,  $\phi_p$  acts as endomorphism of  $E[\ell]$ .

- *division polynomials* of  $E$ : Certain sequence of polynomials, so that

$$(X, Y) \in E[\ell] \Leftrightarrow \psi_\ell(X) = 0$$

holds.

# Bounds for $\#E(\mathbb{F}_p)$

## Theorem (Hasse bound (1933))

Let  $E$  be an elliptic curve over  $\mathbb{F}_p$ . Then

$$p + 1 - 2\sqrt{p} \leq \#E(\mathbb{F}_p) \leq p + 1 + 2\sqrt{p}.$$

Hence  $\#E(\mathbb{F}_p) = p + 1 - t$ , where  $t \in \mathbb{Z}$  and  $|t| \leq 2\sqrt{p}$ .

## Theorem

The Frobenius endomorphism satisfies the quadratic equation

$$\chi(\phi_p) := \phi_p^2 - t\phi_p + p = 0.$$

↪ Schoof's algorithm

# Further considerations

- Consider action of  $\phi_p$  on  $E[\ell]$ .
- Consider roots of  $\chi_\ell(\phi_p) = \phi_p^2 - t\phi_p + p \pmod{\ell}$ .
  - ① Two roots in  $\mathbb{F}_\ell \rightarrow \ell$  is an *Elkies prime*.
  - ② No root in  $\mathbb{F}_\ell \rightarrow \ell$  is an *Atkin prime*.
- In the first case  $\chi_\ell(X)$  has a linear factor over  $\mathbb{F}_\ell[X]$   
 $\rightarrow \psi_\ell(X)$  has factor  $f_\ell(X) = \prod_{a=1}^{(\ell-1)/2} (X - (aP)_x)$  where  
 $\varphi_p(P) = \lambda P$ .  
 $\rightsquigarrow$  Elkies procedure with improved run-time

# Elliptic Gauss sum

Let  $\chi$  be a Dirichlet character of order  $n \mid \ell - 1$ , then we define an *elliptic Gauss sum* (Mihailescu) as

$$\tau_e(\chi) = \sum_{a=1}^{\ell-1} \chi(a)(aP)_v, \quad \begin{cases} v = x, & n \equiv 1 \pmod{2}, \\ v = y, & n \equiv 0 \pmod{2}. \end{cases}$$

## Lemma

The elliptic Gauss sum has the following properties:

- 1  $\tau_e(\chi)^n \in \mathbb{F}_p[\zeta_n]$
- 2  $\varphi_p(\tau_e(\chi)) = \chi^{-p}(\lambda)\tau_e(\chi^p)$



# Modular functions I

## Definition

- 1 Upper half-plane  $\mathbb{H} := \{\tau \in \mathbb{C} : \Im(\tau) > 0\}$ .
- 2  $\Gamma = \mathrm{SL}_2(\mathbb{Z})$  acts on  $\mathbb{H}$  via

$$\gamma = \begin{pmatrix} a & b \\ c & d \end{pmatrix} : \mathbb{H} \rightarrow \mathbb{H}, \quad \tau \mapsto \frac{a\tau + b}{c\tau + d}.$$

# Modular functions II

## Definition

Let  $f(\tau)$  be a meromorphic function on  $\mathbb{H}$ ,  $k \in \mathbb{Z}$ . We call  $f(\tau)$  a *modular function of weight  $k$*  for  $\Gamma' \subseteq \mathrm{SL}_2(\mathbb{Z})$  (where we require  $\begin{pmatrix} 1 & N \\ 0 & 1 \end{pmatrix} \in \Gamma'$  for some  $N \in \mathbb{N}$ ) if it satisfies the following conditions

- 1  $f(\gamma\tau) = (c\tau + d)^k f(\tau) \quad \forall \gamma = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \Gamma'$ .

In particular, this implies there is a Laurent series for  $f(\tau)$  in terms of  $q_N = \exp\left(\frac{2\pi i\tau}{N}\right)$ .

- 2 In the Laurent series for  $f(\gamma\tau) = \sum_{n \in \mathbb{Z}} a_n q_N^n$  we have  $a_n = 0$  for  $n < n_0$ ,  $n_0 \in \mathbb{Z} \quad \forall \gamma \in \mathrm{SL}_2(\mathbb{Z})$ .

In applications we focus on

$$\Gamma' = \Gamma_0(\ell) = \left\{ \gamma = \begin{pmatrix} a & b \\ c & d \end{pmatrix} \in \mathrm{SL}_2(\mathbb{Z}) : c \equiv 0(\ell) \right\}, \quad \ell \text{ prime.}$$

# Examples

$$E_{2k}(\tau) = \frac{1}{\zeta(2k)} \sum'_{n,m \in \mathbb{Z}} \frac{1}{(m + n\tau)^{2k}} \text{ for } k > 1,$$

$$\Delta(\tau) = \frac{E_4(\tau)^3 - E_6(\tau)^2}{1728},$$

$$j(\tau) = \frac{E_4(\tau)^3}{\Delta(\tau)},$$

$$\eta(q) = q^{\frac{1}{24}} \prod_{n=1}^{\infty} (1 - q^n),$$

$$m_\ell(\tau) = \ell^s \frac{\eta(\ell\tau)^{2s}}{\eta(\tau)}, \quad s = \min_{s \in \mathbb{N}} \left\{ s : \frac{s(\ell-1)}{12} \in \mathbb{N} \right\},$$

$$j(\ell\tau).$$

# Facts on modular functions

## Lemma

*Modular functions of weight 0 form a field  $\mathbf{A}_0(\Gamma')$ .*

## Theorem

*With notation as on the last slide, we have*

- 1  $\mathbf{A}_0(\Gamma) = \mathbb{C}(j)$ ,
- 2  $\mathbf{A}_0(\Gamma_0(\ell)) = \mathbb{C}(j, f)$  for  $f \in \mathbf{A}_0(\Gamma_0(\ell)) \setminus \mathbb{C}(j)$ .

So, given  $g \in \mathbf{A}_0(\Gamma_0(\ell))$ , there exist  $P_1, P_2 \in \mathbb{C}[X, Y]$  s. t.

$$g = \frac{P_1(f, j)}{P_2(f, j)}$$

We now focus on  $f = m_\ell(\tau)$ .

## Facts on modular functions II

## Lemma (B)

Let  $g \in \mathbf{A}_0(\Gamma_0(\ell))$  be holomorphic. Then  $g$  admits a representation of the form

$$g(\tau) = \frac{Q(m_\ell, j)}{m_\ell^k \frac{\partial G_\ell}{\partial Y}(m_\ell, j)},$$

for some  $k \geq 0$  and a polynomial  $Q(X, Y) \in \mathbb{C}[X, Y]$ , where

$$\deg_Y(Q) < \deg_Y(G_\ell) = \min_{s \in \mathbb{N}} \left\{ v = \frac{s(\ell - 1)}{12} : v \in \mathbb{N} \right\}$$

and  $G_\ell(X, j)$  is the minimal polynomial of  $m_\ell$  over  $\mathbb{C}(j)$ .

# Tate curve

## Proposition (Tate)

Let  $E_4, E_6$  be as before. Then the quantities

$$x(w, q) = \frac{1}{12} + \frac{w}{(1-w)^2} + \sum_{n=1}^{\infty} \sum_{m=1}^{\infty} mq^{nm}(w^m + w^{-m}) - 2mq^{nm},$$

$$y(w, q) = \frac{w + w^2}{2(1-w)^3} + \frac{1}{2} \sum_{n=1}^{\infty} \sum_{m=1}^{\infty} \frac{m(m+1)}{2} (q^{nm}(w^m - w^{-m}) + q^{n(m+1)}(w^{m+1} - w^{-(m+1)}))$$

satisfy

$$E_q : \quad y(w, q)^2 = x(w, q)^3 - \frac{E_4(q)}{48}x(w, q) + \frac{E_6(q)}{864}.$$

$E_q$  is called the Tate curve which parametrizes isomorphism classes of elliptic curves over  $\mathbb{C}$ .

# Universal elliptic Gauss sums

## Lemma (B)

Let  $\ell$  be a prime,  $n \mid \ell - 1$ ,  $\chi : \mathbb{F}_\ell^* \mapsto \mu_n$  a Dirichlet character,  $\zeta$  an  $\ell$ -th root of unity and let  $r, e_\Delta$  be appropriately chosen integers. Let in addition  $V = x$  for odd and  $V = y$  for even  $n$  and define

$$G_{\ell,n}(q) = \sum_{\lambda \in \mathbb{F}_\ell^*} \chi(\lambda) V(\zeta^\lambda, q), \quad p_1(q) = \sum_{\lambda \in \mathbb{F}_\ell^*} x(\zeta^\lambda, q).$$

Then

$$\tau_{\ell,n}(q) := \frac{G_{\ell,n}(q)^n p_1(q)^r}{\Delta(q)^{e_\Delta}},$$

is a modular function of weight 0 for  $\Gamma_0(\ell)$ , holomorphic on  $\mathbb{H}$  and has coefficients in  $\mathbb{Q}[\zeta_n]$ . We call it a universal elliptic Gauss sum.

## Proof.

Study behaviour of Weierstraß  $\wp$ -function under action of  $\mathrm{SL}_2(\mathbb{Z})$  and use connection between  $x(w, q), y(w, q)$  and  $\wp(z, \tau), \wp'(z, \tau)$ . □

# An algorithm for computing

By general lemma we find

$$\tau_{\ell,n}(q) = \frac{Q(m_\ell, j)}{m_\ell^k \frac{\partial G_\ell}{\partial Y}(m_\ell, j)}.$$

So use the following algorithm:

- 1 Compute  $\tau_{\ell,n}(q) \frac{\partial G_\ell}{\partial Y}(m_\ell, j) =: s$  up to precision  $\text{prec}(\ell, n)$ ,  $Q := 0$ .
- 2 Determine  $o = \text{ord}(s)$  and  $(i, k) : iv - k = o$  and  $k < v$ .
- 3 Compute  $s := s - cm_\ell^i j^k$ ,  $Q := Q + cX^i Y^k$
- 4 Repeat 2 and 3 until  $s = 0$ .



# Required precision

## Lemma (B.)

*We can take  $\text{prec}(\ell, n) = (v + e_\Delta)\ell$ .*

Run-time:

- Compute  $\tau_{\ell,n}(q)$ :  $\tilde{O}(\ell n v)$
- Determine  $Q$ :  $\tilde{O}(\ell^2 v^2)$

# Application

Recall Schoof's algorithm (1985)

- Compute  $\#E(\mathbb{F}_p) = p + 1 - t$ ,  $|t| \leq 2\sqrt{p}$
  - $\rightsquigarrow$  Determine  $t \pmod{\ell}$  for small primes  $\ell$  by finding  $t$  s. t.  $\varphi_p^2 - t\varphi_p + p \equiv 0 \pmod{\ell}$ , then use CRT
  - First polynomial algorithm (in  $\log p$ )
  - If  $\ell$  is Elkies prime: Use polynomials of lower degree  $\Rightarrow$  power saving in run-time
  - If  $\ell$  is Atkin prime: Generic approach of equal run-time + sophisticated BSGS
- $\rightsquigarrow$  SEA combines Elkies (mostly) + Atkin procedures

# Elkies procedure

Need to find  $\lambda$  s. t.  $\varphi_p(P) = \lambda P$  for  $\ell$ -torsion point  $P$ .

$\rightsquigarrow$  Compute in  $\mathbb{F}_p[X]/(f_\ell(X))$ , extension of degree  $\mathcal{O}(\ell)$ .

## Lemma (Mihalescu, 2006)

Let  $\ell$  be a prime,  $\chi$  be a character with  $\text{ord}(\chi) = n \parallel \ell - 1$ . Let  $\tau_e(\chi)$  be the elliptic Gauss sum. Then

$$\varphi_p(\tau_e(\chi)) = \chi^{-p}(\lambda) \tau_e(\chi^p)$$

Writing  $p = nq + m$ , one obtains

$$(\tau_e(\chi))^n)^q \cdot \frac{\tau_e(\chi)^m}{\tau_e(\chi^m)} = \chi^{-m}(\lambda)$$

Both factors lie in  $\mathbb{F}_p[\zeta_n]$ , computations can be done in extension of degree  $\mathcal{O}(n)$  and no searching for  $\lambda$  is required.

# Compute the factors

Use universal elliptic Gauss sums: We know

$$\tau_{\ell,n}(q) = \frac{G_{\ell,n}(q)^n p_1(q)^r}{\Delta(q)^{e_\Delta}} = R(j(q), m_\ell(q)).$$

Substitute  $q = \exp(2\pi i\tau(E)) \Rightarrow \tau_{\ell,n}(E) = R(j(E), m_\ell(E))$  for curve  $E$  in question.

Hence, compute  $j(E), \Delta(E), p_1(E)$  and obtain  $m_\ell(E)$  as root of  $G_\ell(X, j(E))$ .  $\rightsquigarrow$  Compute  $\tau_e(\chi)^n$  for our  $E$

Similar approach for Jacobi sums  $\rightsquigarrow$  determine  $\lambda \pmod n$  for all  $n \mid \ell - 1$ .  
CRT gives index of  $\lambda$  in  $(\mathbb{Z}/\ell\mathbb{Z})^*$   $\rightsquigarrow t = \lambda + p/\lambda \pmod \ell$  and  $t$ .

# Further research

- 1 Replace  $m_\ell$  by other modular functions to improve run-time
- 2 Analyse coefficient size
- 3 ?

Merci pour votre attention