

Finding optimal Chudnovsky-Chudnovsky multiplication algorithms

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A trick

$$\begin{aligned} & \text{Compute } (ax + b)(cx + d) \\ &= a \bullet cx^2 + (a \bullet d + b \bullet c)x + b \bullet d \end{aligned}$$

Total : **4** multiplications

Can one do
better ?

Compute $(ax + b)(cx + d)$

$$= a \bullet cx^2 + (a \bullet d + b \bullet c)x + b \bullet d$$

Total : **4** multiplications

A trick

Evaluate $m_0 = b \bullet d$

$$m_1 = (a + b) \bullet (c + d)$$

$$m_\infty = a \bullet c$$

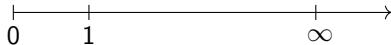
Then, $(ax + b)(cx + d) = m_\infty x^2 + (m_1 - m_0 - m_\infty)x + m_0$

Total : **3** multiplications

What happened ?

Lagrange's interpolation (over $\mathbf{R} \cup \infty$)

The degree 2 polynomial $P(x)=(ax+b)(cx+d)$ is fully determined by the 3 evaluations $m_0 = P(0)$, $m_1 = P(1)$, $m_\infty = P(\infty)$.



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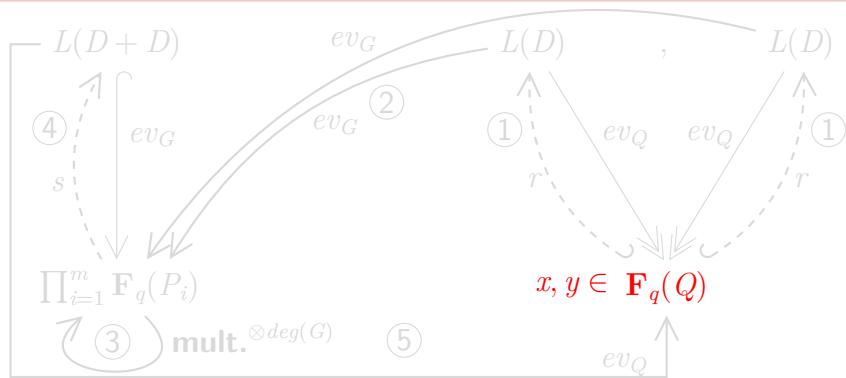
**Problem : only $q + 1$
points on $\mathbf{F}_q \cup \infty$**



Ch&Ch's interpretation

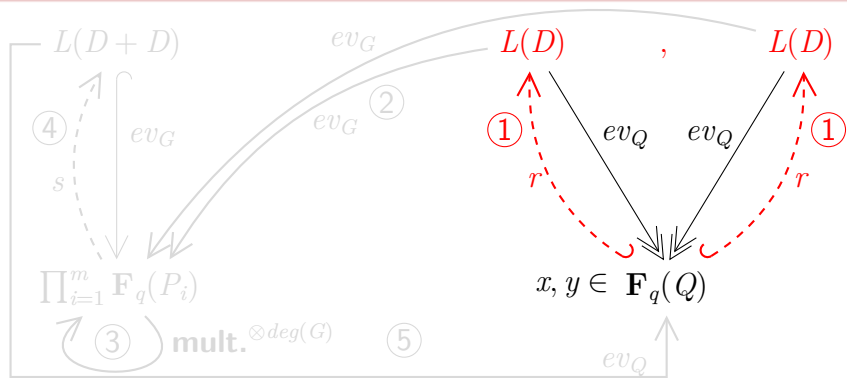
	Before	After
set:	$\mathbf{F}_q \cup \infty$	curve $X = P_{\mathbf{F}_q}^1$
$(ax + b)$ and $(cx + d)$:	polynomials	sections of $D = \mathcal{O}_X(\infty)$
evaluation on:	points $0, 1, \infty$	divisor $G =$ $[0] + [1] + [\infty]$

Multiply x, y in \mathbf{F}_{q^m} (Ch&Ch)



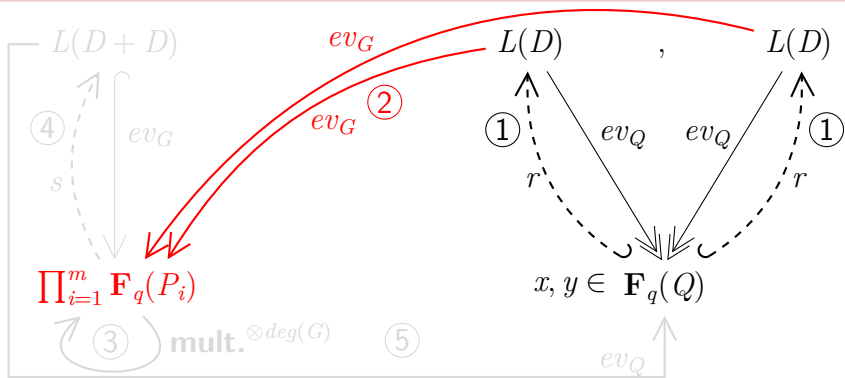
① choose Q on X of degree m , fix isomorphism $x, y \in \mathbf{F}_{q^m} \cong \mathbf{F}_q(Q)$

Multiply x, y in \mathbf{F}_{q^m} (Ch&Ch)



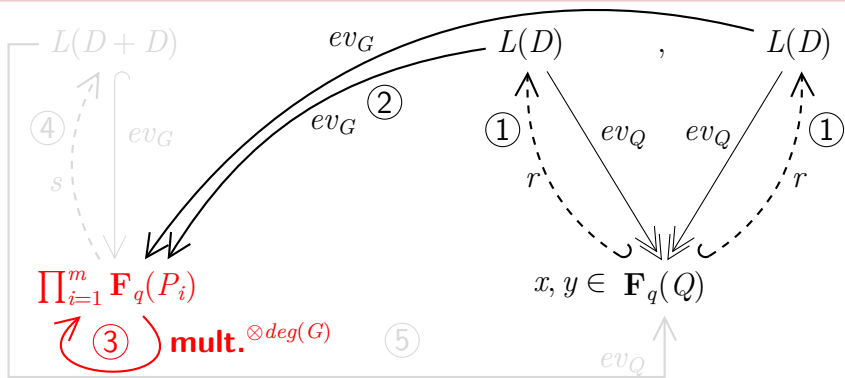
- ① choose Q on X of degree m , fix isomorphism $x, y \in \mathbf{F}_{q^m} \cong \mathbf{F}_q(Q)$
- ① find divisor D ; lift x, y to f_x, f_y in $L(D)$.

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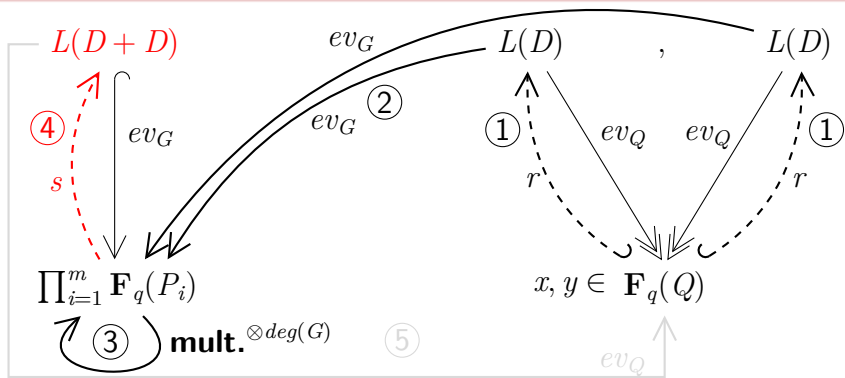
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- ② find divisor $G = P_1 + \dots + P_n$; evaluate the $f_x(P_i)$ and $f_y(P_i)$.

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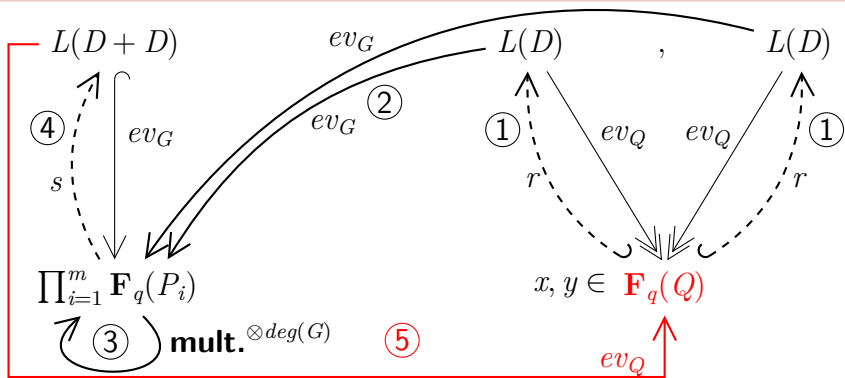
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- 1** find divisor D ; lift x, y to f_x, f_y in $L(D)$.
- 2** find divisor $G = P_1 + \dots + P_n$; evaluate the $f_x(P_i)$ and $f_y(P_i)$.
- 3** compute the $m_i = f_x(P_i) \bullet f_y(P_i)$: **deg G** multiplications.

Multiply x, y in \mathbf{F}_{q^m} (Ch&Ch)



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- ⑤ interpolate $[m_1, \dots, m_n]$ to unique $g \in L(D+D)$
- ⑥ evaluate g at Q to find the product of x and y .

Need more interpolation data ?

	Before	After	Evaluation in:
degree(P_i) :	1	$d \geq 1$	\mathbf{F}_{q^d}
" $f_x(P_i)$ " :	value $f_x(P_i)$	derivatives of f_x at P_i up to $l \geq 0$	$\frac{\mathbf{F}_q[y]}{y^l}$
...Both :			$\frac{\mathbf{F}_{q^d}[y]}{y^l}$

Putting things together

Note $\mu_q^{\text{sym}}(d, l)$ the bilinear complexity of the multiplication in $\frac{\mathbf{F}_{q^d}[y]}{y^l}$.

Complexity of the algorithm [Randriam 2012, see Th. 2]

X a curve of genus g over \mathbf{F}_q , Q a point of degree m , D a divisor, $G := l_1 P_1 + \dots + l_n P_n$ [with $\deg P_i = d_i$], and **suppose** (G, D, Q) **"suitable for interpolation"**. Then :

$$\mu_q^{\text{sym}}(m) \leq \sum_{i=1}^n \mu_q^{\text{sym}}(d_i, l_i)$$

weighted degree of G

Find suitable (G, D, Q) , G of
smallest weighted degree



(G, D, Q) , G of smallest degree?

Best expectable (G, D, Q) [see R., Prop. 8 & 10]

X a curve of genus g . Assume $m > g$. Then : criterion for (G, D, Q) being "suitable for interpolation on \mathbb{F}_{q^m} " depends only on the **classes** of G, D, Q in $\text{Cl}(X)$. When the case :

1. $\deg G \geq 2m + g - 1$
2. ...and if this lower bound attained, then $\deg D = m + g - 1$.

Example [see 4.2]

Multiplication in \mathbf{F}_{2^m} , $m = 163$:

Step 0 : find a curve X having a divisor G of **smallest weighted degree, under the constraint $\deg G \geq 2m + g - 1 = 331$** (condition 1.).

→ Exhaustive search on the $X_0(N)$...Winner : $X_0(71)$, with a G of weighted degree 900.



G is not necessarily part of a suitable (G, D, Q)

Example [see 4.2]

class group : $\text{Cl}(X_0(71)) \sim \mathbf{Z}/315\mathbf{Z} \times \mathbf{Z}$, generators : (D_1, D_2) .

Step 1 : Choose a G on X of weighted degree 900.

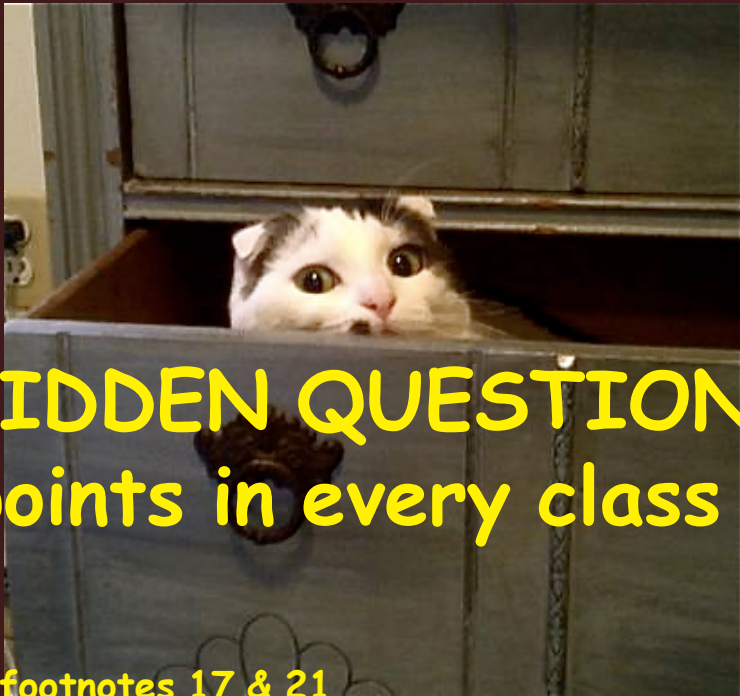
- Step 2 $\deg D$ must be $m + g - 1 = 168 \rightarrow$ Loop over the classes :
 $D := i * D_1 + 168 * D_2$, until $l(2D - G) = 0$ [Theorem 2, condition (i')]. \rightarrow Success for $i = 2$.
- Step 3 : Loop over the classes of random Q of degree m , until $l(D - Q) = 0$ [Theorem 2, condition (ii') & Footnote 16] \rightarrow Success at first attempt.

\rightarrow 900 is a new upper bound for the multiplication in $\mathbf{F}_{2^{163}}$.

Search on the curves $X_0(N)$, $N = 0 \dots 1000 \longrightarrow$ new bounds:

Table: New upper bounds on $\mu_2^{\text{sym}}(m)$, sorted by the genus of curves used

$m \backslash g$	1 (ECs)	2	3	4	5	6
163	905	903	901	.	.	900
233	1339	1336	.	1335	.	.
283	1661	1660	.	1654	.	.
409	2492	2491	.	2486	.	.
571	3562	3561	3560	3555	.	.



HIDDEN QUESTION :
points in every class ?

...see footnotes 17 & 21